

An Incentive Mechanism for Message Relaying in Peer-to- Peer Discovery

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Peer-to-Peer Discovery

- Peer discovery is important in a P2P system for peers to exploit distributed resources
- Message relaying for peer discovery in a distributive model (“word-of-mouth”)
- Common searching protocols:
 - breadth-first search (Gnutella)
 - depth-first search (Freenet)

Problems

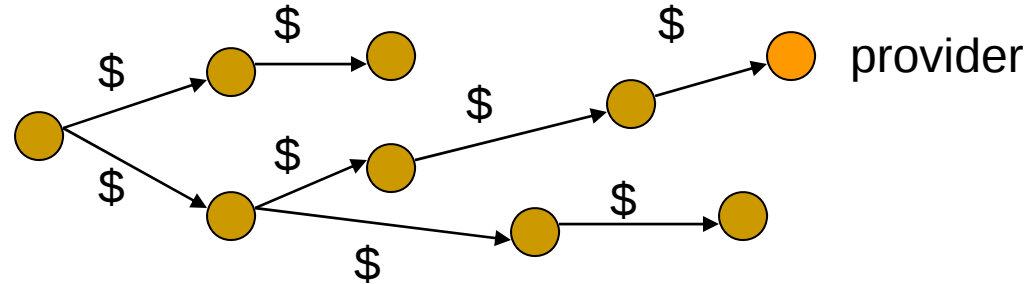
■ Efficiency

- ❑ BFS costs enormously the bandwidth although the results can be found quickly
- ❑ DFS is cheaper in communications but the response time can be very long

■ Incentive

- ❑ Each peer may represent a self-interested entity
- ❑ Communication bandwidth and energy are bounded resources
- ❑ A peer may drop searching messages from other peers to save resources (*free riding*)
- ❑ A *different* problem from the message relaying mechanism in distributed *routing*

The Incentive Mechanism



- An incentive mechanism overcomes the flooding problem of BFS while reserving the quick response feature
- It is motivated by the following requirements:
 - Communication efficiency
 - Reliability
 - Anonymity and information locality

Communication Efficiency

- *Overlapping*: the wasted transmissions of the message to the peers that have received the message
- The system becomes *saturated* quickly if each peer makes significant transmission efforts
- To reduce the inefficiency caused by overlapping, a peer should adjust the effort with the saturation status.

Reliability

- The probability of finding a provider is positively correlated with the transmission efforts.
- Reliability is a conflicting goal with reducing the communication cost.
- Trade off the communication cost and reliability to maximize the utility.
- The optimum depends on the value of information.

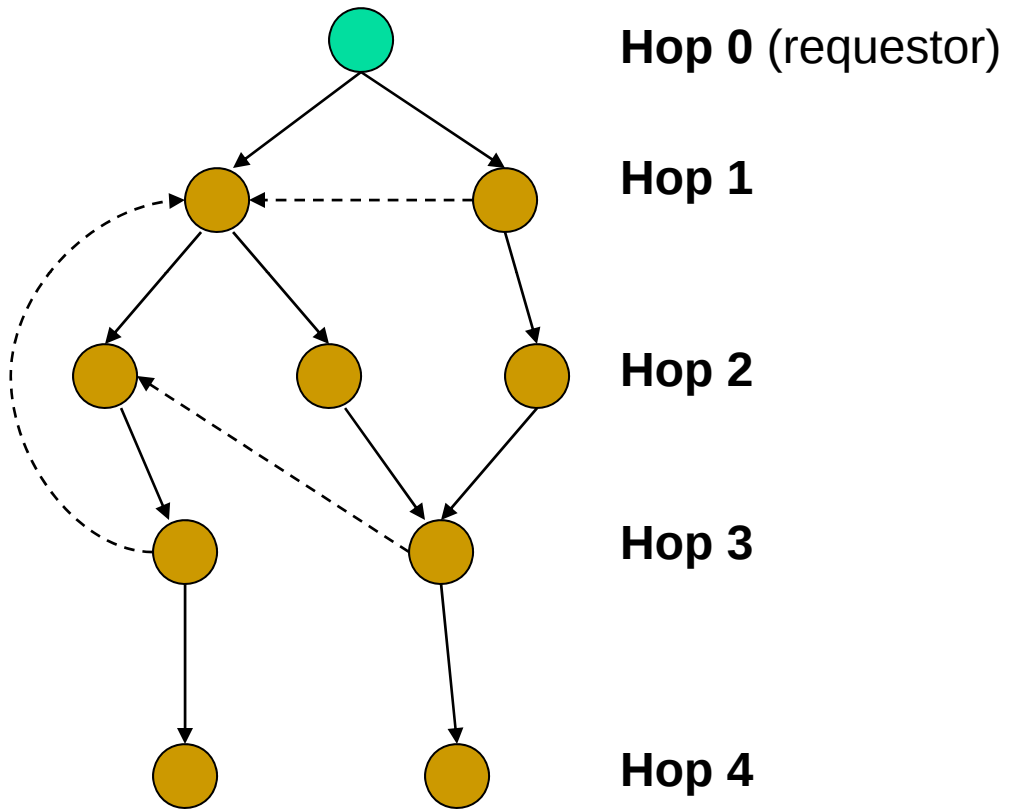
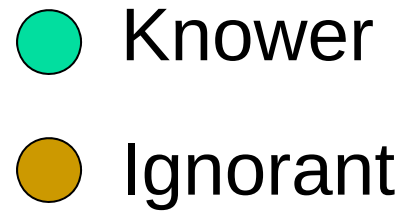
Anonymity and Information Locality

- A micro-payment system that prices the scarce resource is commonly used to provide incentive compatibility.
- The usual micro-payment mechanism is NOT feasible in P2P discovery system:
 - With *anonymity* the intermediate peers on the route and their transmission efforts are not identifiable by the source node.
 - Pricing is difficult as the *local* environments are unknown to the mechanism designer or a third party (“non-private value revelation” in Shneidman & Parkes).
- Price the searching result that only require local information that is easy to obtain/verify by local price makers, e.g., neighbors, immediate upstream nodes, and incentives received and passed.

Agenda

- The relaying model
- Equilibrium analysis
- Experiments
- Conclusion

The Relaying Model



The Relaying Model (cont.)

■ Assumptions

- ❑ Single provider
- ❑ Homogeneity

■ Parameters

- ❑ Message relaying cost: $c(k)$, increasing and convex
- ❑ Number of peers: N
- ❑ Max hop number (TTL): H , i.e., $H+1$ hops allowed
- ❑ Information value: v_0

Equilibrium Analysis

- Individual strategy:

$$(k_i, u_i) = S_i(h_i, v_i)$$

- Individual utility:

$$U_i(S) = (v_i - u_i) L_i(S) / N - c(k_i)$$

- System utility:

$$U(S) = v_0 L_0(S) / N - \sum c(k_i)$$

- Sub-game perfect Nash Equilibrium (SPNE): a Nash equilibrium for each subgame of the propagation process starting from each hop to the last hop.

Equilibrium Analysis (cont.)

PROPOSITION 2.1. ⁴ *Given the incentive relaying strategy $u_i(h_i, v_i)$ for each peer i , a SPNE of the transmission strategies $\{k_i^*(h_i, v_i|u_i)\}_{i \in N}$ exists. $k_i^*(h_i, v_i|u_i)$ decreases with the increase of h_i or u_i , or with the decrease of v_i .*

PROPOSITION 2.2. *With a best response strategy the increase of the input incentive leads to the increase of the expected number of downstream nodes.*

PROPOSITION 2.3. *If all peers have the same degree, a symmetric SPNE exists and is unique.*

Approximation of Symmetric SPNE

- The expected ignorants reached by each peer in hop h is

$$\hat{k}_h = \begin{cases} 0 & \text{if } k_h = 0 \text{ or } m_h = 0 \\ \frac{N - n_h}{m_h} [1 - (1 - \frac{k_h}{N-1})^{m_h}] & \text{else.} \end{cases}$$

- Estimate the number of peers in hop $h+1$:

$$m_{h+1} \doteq m_h \hat{k}_h = (N - n_{h-1} - m_h) [1 - (1 - \frac{k_h}{N-1})^{m_h}]$$

- Estimate the number of descendants of a peer in hop h :

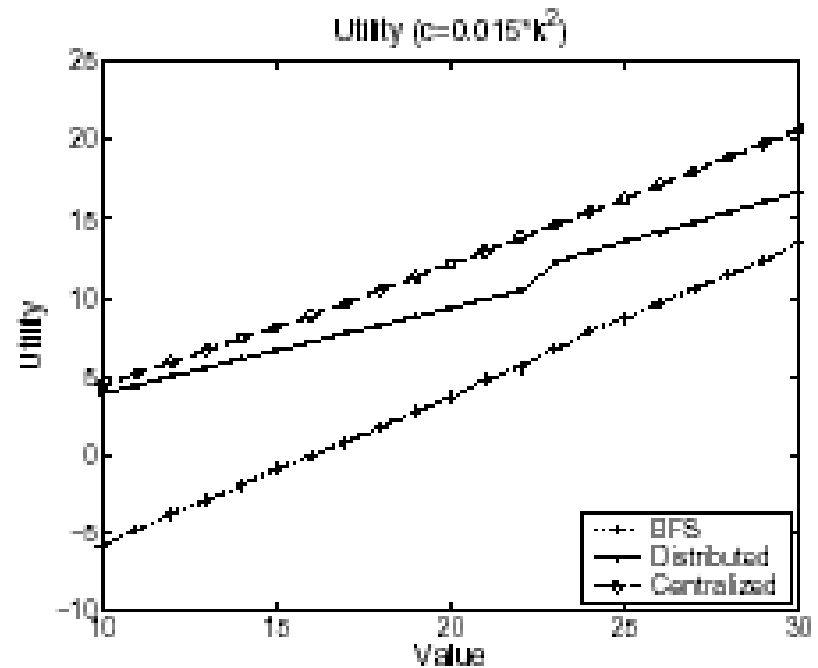
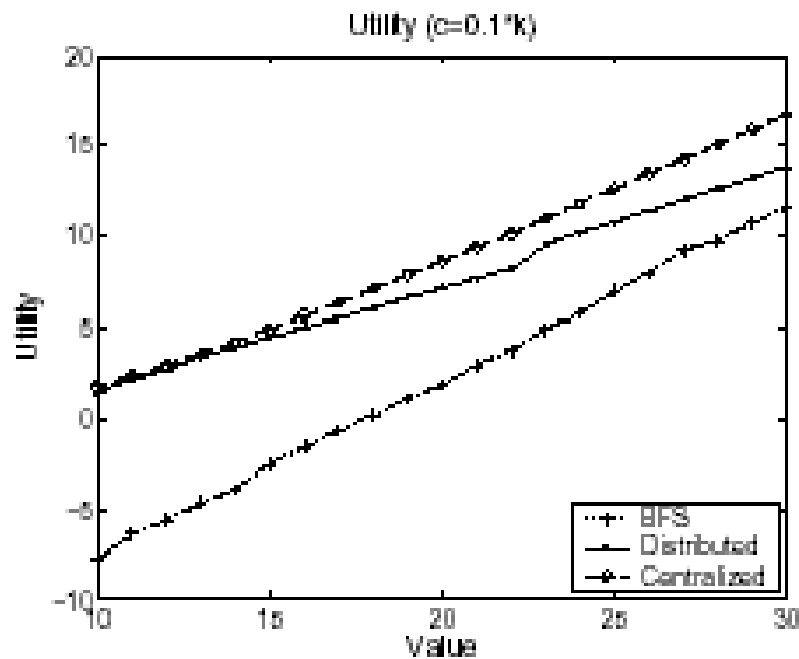
$$L_h(n_h, m_h) = \sum_{l=h+1}^H m_l / m_h$$

- An approximate SPNE can be calculated by backward induction.

Experiments

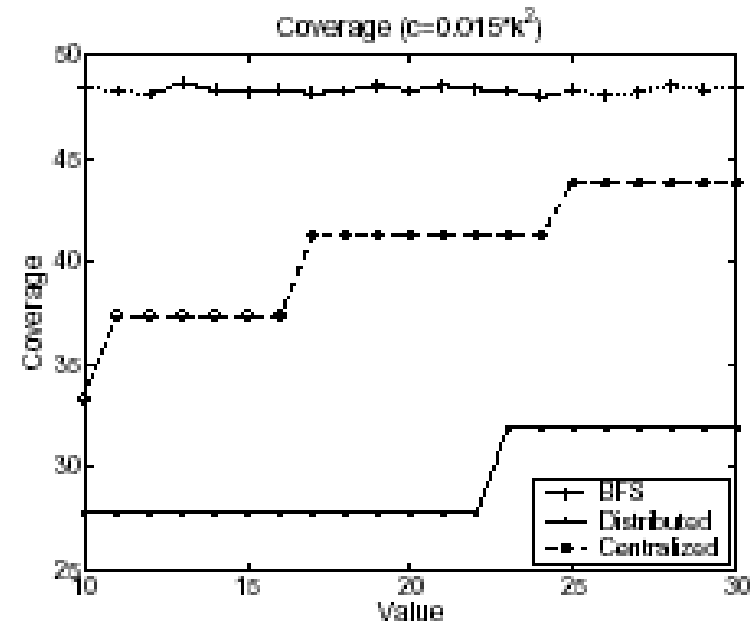
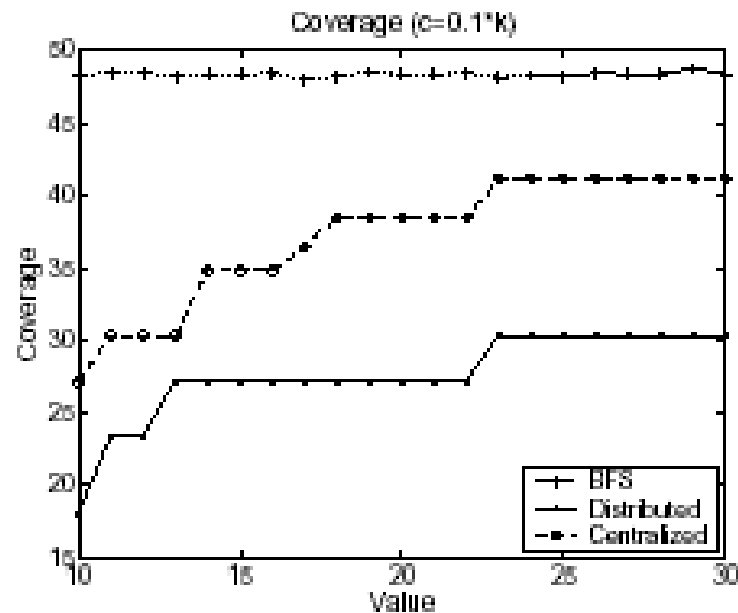
- Three systems:
 - The distributed incentive mechanism (based on the approximate SPNE strategy)
 - The breadth-first search mechanism
 - The centralized mechanism (maximizes the system utility based on the approximate coverage function)
- We are interested in:
 - The system total utility (efficiency)
 - The coverage (reliability)
 - The distribution of transmission effort and incentives over hops in the distributed incentive mechanism
- $N=50$, $H=2$ (three hops), $D=6$, $v_0=10$ to 30 , $c(k)=0.1k$ or $0.015k^2$

System Utility



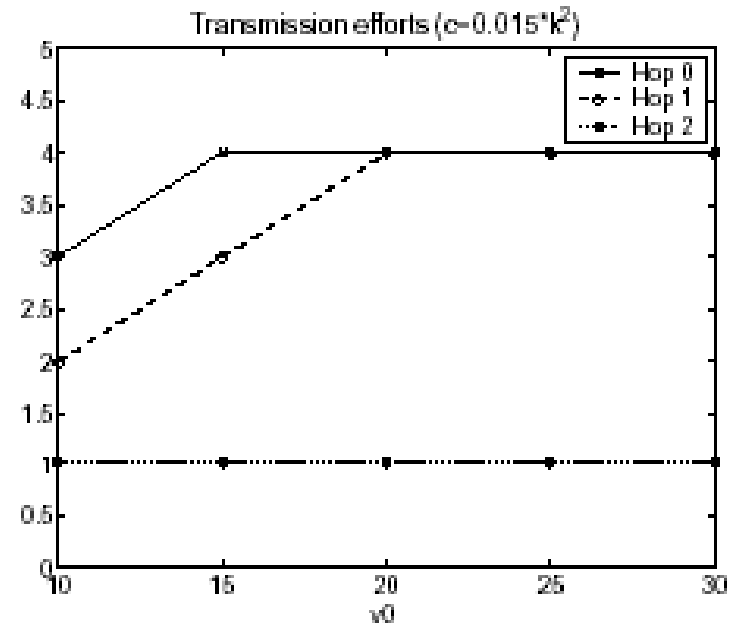
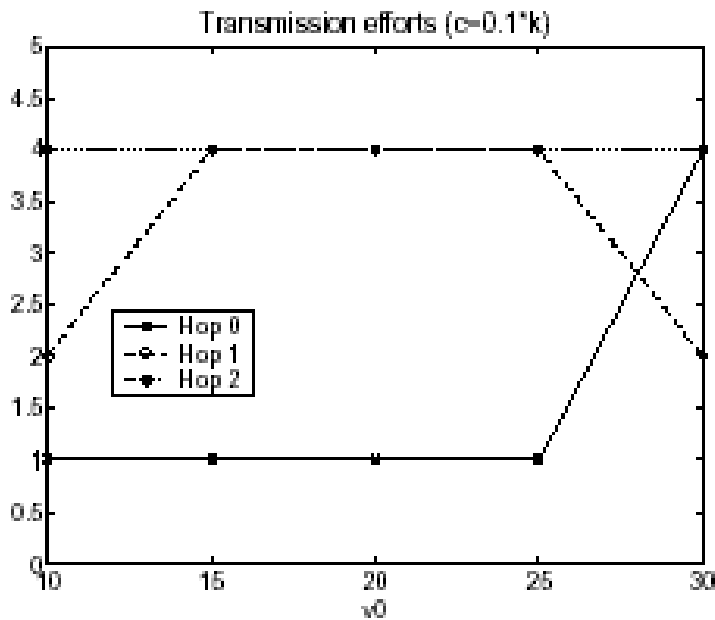
- $U_{BFS} < U_{dist} < U_{cen}$ with U_{dist} generally higher than 80% of U_{cen}
- $U_{dist} - U_{BFS}$ decreases with the increase of information value.

Coverage



- BFS has the highest coverage, which does not depend on the information value.
- The coverage of both the distributed and centralized system increases with the information value
- The distributed mechanism cannot achieve the optimal coordination.

Distribution of Transmission Efforts and Incentives



With a more convex cost function, a peer tends to develop its family tree by its own transmission efforts, and pass little incentives to downstream nodes.

Conclusion

- An incentive mechanism for message relaying in P2P discovery that prices the searching *results* instead of the searching *behavior*.
- Optimal TTL?
- Multiple providers?
- Peers are heterogeneous and have knowledge?